Cache-aware response time analysis for real-time tasks with fixed preemption points

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Abstract—In real-time systems that employ preemptive scheduling and cache architecture, it is essential to account as precisely as possible for cache-related preemption delays in the schedulability analysis, as an imprecise estimation may falsely deem the system unschedulable. In the current state of the art for preemptive scheduling of tasks with fixed preemption points, the existing schedulability analysis considers overly pessimistic estimation of cache-related preemption delay, which eventually leads to overly pessimistic schedulability results. In this paper, we propose a novel response time analysis for real-time tasks with fixed preemption points, accounting for a more precise estimation of cache-related preemption delays. The evaluation shows that the proposed analysis significantly dominates the existing approach by being able to always identify more schedulable tasksets.

Index Terms—Real-time Systems, Cache-related preemption delay (CRPD), Fixed-Priority scheduling, Preemptive Scheduling, Cache Memory

I. INTRODUCTION

For many years it has been debated whether the non-preemptive real-time scheduling or the fully-preemptive real-time scheduling offers better schedulability bound and, if so, under which circumstances. Eventually, it was shown by Yao et al. [1], and Butazzo et al. [2] that Limited-Preemptive Scheduling (LPS) can provide even better schedulability results than the two since it can combine and mimic both of the scheduling paradigms.

Under the LPS paradigm, several approaches emerged, such as Deferred Preemption Scheduling (LP-DPS) proposed by Sanjoy Baruah [3], Preemption Thresholds Scheduling (LP-PTS) proposed by Wang and Saksena [4], and Fixed Preemption Points (LP-FPP) Scheduling, proposed by Alan Burns [5]. Among all of them, it has been shown by Yao et al. [1], and Butazzo et al. [2] that LP-FPP allows for the most precise estimation of the taskset schedulability since the possibly preempted points of a task are limited and known in advance, prior the system’s runtime.

It is shown by Pellizoni et al. [6] that preemptions can introduce a significant preemption related delay in preemptive real-time systems, even up to 33% of the task’s worst-case execution time. Also, it is shown by Bastoni et al. [7] that in real-time systems which employ a cache-based architecture, cache-related preemption delay (CRPD) is the major part of the preemption related delay. Those are a few reasons why it can be important to limit the number of preemptions, which is possible by using LP-FPP scheduling, and why such systems may benefit from a precise CRPD estimation since the preemption still may occur and a pessimistic estimation may deem a schedulable taskset being unschedulable.

The existing work on cache-related preemption estimation, integrated with the schedulability analysis, has, so far, considered only fully-preemptive real-time systems, and has been shown by Altmeyer et al. [8], [9] to be highly beneficial in determining the schedulability of real-time systems. However, concerning the LP-FPP scheduling, the existing feasibility analysis proposed by Yao et al. [10], [11] accounts for overly pessimistic estimation of preemption related delay, without seizing many of the properties of the LP-FPP task model which facilitate the CRPD estimation of a task.

In this paper, we propose a novel cache-aware response time analysis for sporadic fixed-priority real-time tasks with fixed preemption points and sequential non-preemptive regions. We first identify reloadable cache blocks, by defining the maximum number of times each memory cache block can be reloaded upon all possible preemptions, which is an improvement compared to the existing methods. Secondly, we define the CRPD bounds accounting for the individual preemptions on the non-preemptive regions, finally combining the two approaches in order to define a safely upper-bounded CRPD estimation which is integrated into the schedulability analysis. Also, we identify and propose a computation of the maximum time interval during which any of the released preemting jobs may affect the CRPD of a task. Evaluation results show that the proposed analysis outperforms the existing feasibility analysis for the LP-FPP task model in all of the evaluated cases, regardless of the variation in different task and cache parameters.

In the remainder of the paper, we first define the system model in Section II. In Section III we discuss the related work and describe the existing feasibility analysis for LP-FPP scheduling. The proposed response time analysis is motivated in Section IV and defined in Section V. The conclusions are discussed in Section VII.
In this paper, we consider a sporadic task model, under a preemptive scheduler. In more detail, a taskset \( \Gamma \) consists of \( n \) tasks with preassigned fixed and disjunct priorities, where each task \( \tau_i \) generates an infinite number of jobs, characterised by the following tuples of parameters: \( \langle P_i, C_i, T_i, D_i \rangle \). For a task \( \tau_i \) (\( 1 \leq i \leq n \)), its priority is denoted with \( P_i \), its worst-case execution time (WCET) without considering for preemptions is denoted with \( C_i \). The minimum inter-arrival time between the consecutive task jobs is \( T_i \), and the relative deadline is denoted with \( D_i \). The assumed deadlines are constrained, i.e. \( D_i \leq T_i \). A j-th job of \( \tau_i \) is denoted with \( \tau_{i,j} \).

We also consider that a task consists of a sequence of \( l_i \) non-preemptive regions (NPR), separated by \( l_i - 1 \) preemption points (see Fig. 1). Non-preemptive regions are also called runnables in some real-time system domains, e.g. the automotive industry [12], [13], [14]. The k-th NPR of \( \tau_i \) is denoted with \( \delta_{i,k} \), \( 1 \leq k \leq l_i \), and its worst-case execution time, assuming no preemptions during the execution of \( \delta_{i,1}, \ldots, \delta_{i,k-1} \), is denoted with \( q_{i,k} \). Thus \( C_i \) can be expressed as \( \sum_{k=1}^{l_i} q_{i,k} \), and the worst-case execution time of the first \( l_i - 1 \) NPRS is denoted with \( E_i \) and is equal to \( E_i = \sum_{k=1}^{l_i-1} q_{i,k} \). The k-th preemption point of \( \tau_i \) is denoted with \( PP_{i,k} \).

For each task \( \tau_i \) we define a set \( ECB_{i,k} \) of evicting cache blocks such that \( ECB_{i,k} \), with \( \tau_i \), its useful memory block, building on the formulation of this paper, we formally represent cache-sets as integers. The adjusted deadlines for using the proposed method in case of level direct-mapped cache, while in Section V-D we enlist the feasible preemption points and preemption patterns in fully-preemptive scheduling. Prior to this work, a very important contribution was made by Bril et al. [23] who for the first time identified the scheduling anomaly known as self-pushing phenomenon (also described later).
The benefits of using a fixed preemption model is described by Butazzo et al. [2], among which are the increased timing predictability and the facilitation of the mutual exclusion problems. Butazzo et al. [2] and Yao et al. [1] showed that LP-FPPS can in the majority of the cases outperform many of the existing scheduling algorithms. Becker et al. [13] showed the benefits of limited-preemptiveness in automotive systems.

Other works in the area of fixed preemptions based task models included the preemption point selection algorithms, proposed by Bertogna et al. [24], Peng et al. [25], and Cavicchio et al. [26]. Cavicchio et al. [26] identified for the first time that a detailed view on CRPD is important for the tasks with fixed preemption points, especially a possibility of cache block to be reloaded. Also, in our previous work [27], [28] considering the tasks with fixed preemptions, we proposed the constraint satisfaction model for analysing CRPD at a task level. However, the existing feasibility analysis proposed by Yao et al. [10], [11] accounts for the overly pessimistic preemption delay estimation, without seizing many of the properties of the LP-FPP task model which may facilitate the CRPD estimation of a task.

In this paper, we propose a novel cache-aware schedulability analysis for fixed-priority preemptive scheduling for sporadic tasks with fixed preemption points. In the following subsection, we first describe the existing feasibility methods.

A. Overview of the feasibility analysis for tasks with non-preemptive regions

The seminal feasibility analysis for tasks with fixed preemption points was proposed by Yao et al. [10], [11], and here we briefly describe their work using Butazzo’s formulations [29].

1) Self-Pushing phenomenon: A self-pushing phenomenon is a scheduling anomaly which may occur in scheduling tasks with non-preemptive regions, identified and proved by Bril et al. [23]. It states that the worst-case response time of a task does not necessarily occur in the first job after the critical instant, but it may occur in one of the succeeding jobs, during the Level-i active period. As proposed by Bril et al. [30]:

Definition 1: A Level-i pending workload at time point \( t \) is the amount of unfinished execution of jobs with a priority higher than or equal to \( P_i \), released strictly before \( t \).

Definition 2: A Level-i active period is an interval \([a, b]\) such that Level-i pending workload is positive for all \( t \in (a, b) \) and zero in \( a \) and \( b \).

Theorem 1: A maximum response time of task \( \tau_i \) under sporadic LP-FPPS is assumed in a level-i active period that is started when \( \tau_i \) has a simultaneous release with all higher priority tasks, and a subjob that produces maximum lower priority blocking starts an infinitesimal time before the simultaneous release. (proved and proposed by Bril et al. [30])

2) Feasibility analysis: For a single job \( \tau_{i,j} \) of \( \tau_i \), feasibility analysis is divided into two parts. First, it is necessary to compute the latest start time \( S_{i,j} \) of the last non-preemptive region. Then, the latest finish time \( F_{i,j} \) of a job is computed. This is performed in two steps because a task can experience the interference from the higher priority tasks prior to the last NPR. But, when the last NPR starts to execute, it cannot be further interfered. The latest start time \( S_{i,j} \) of the last NPR of \( \tau_{i,j} \) is equal to the least fixed point of the following equation:

\[
\begin{align*}
S^{(0)}_{i,j} &= b_i + C_i - q_i I_i + \sum_{\forall r_k \in \text{lp}(i)} C_h \\
S^{(r)}_{i,j} &= b_i + j \times C_i - q_i I_i + \sum_{\forall r_k \in \text{lp}(i)} \left( \frac{S^{(r-1)}_{i,j}}{T_h} \right) + 1 \times C_h
\end{align*}
\]

where \( b_i \) denotes the maximum lower priority blocking that can be imposed on \( \tau_i \). The latest finish time \( F_{i,j} \) of the j-th job is computed with the following equation \( F_{i,j} = S_{i,j} + q_i I_i \), i.e. by adding the latest start time of the j-th job’s last NPR and the worst-case execution time \( q_i I_i \) of the NPR \( \delta_i I_i \).

The worst-case response time \( R_i \) is finally defined as the maximum latest finish time of all jobs released during the Level-i active period.

Furthermore, Yao et al. [11] proved that the analysis can be simplified to the first job of each task, if for each task \( D_i \leq T_i \), and if the taskset is feasible under fully-preemptive scheduling.

B. Limitations of the existing approach

Considering the preemption related overhead, Yao et al. [11] proposed that all \( C_i \) values should be updated with the following equation: \( C_i = C_i^{pp} + (l_i - 1) \times \epsilon_i \), where \( C_i^{pp} \) is the worst-case execution time without considering preemptions, and \( \epsilon_i \) is the largest possible preemption cost that \( \tau_i \) can experience at one of its preemption points. This can be a significant overapproximation because:

\( \diamond \) the analysis omits the fact that CRPD of a task \( \tau_i \) can be caused only upon the execution start of the first NPR \( \delta_i I_i \), until the latest start time of the last NPR \( \delta_i I_i \).

\( \diamond \) the analysis does not account for the case where a preempting task \( \tau_h \) can preempt only a limited number of preemption points of a preempted task \( \tau_i \). Instead, it implicitly assumes that each preempting task can preempt \( \tau_i \) at all of its preemption points.

\( \diamond \) the analysis does not account for the case that between the two accesses of a memory block, it can be reloaded at most once due to preemption.

Therefore, in the following section, we provide motivating examples for the novel method and then we define a response time analysis with a more precise CRPD estimation.

IV. MOTIVATING EXAMPLES

In this section, we present two motivating examples. In Fig. 2, we show a high-level scheduling perspective of a job \( \tau_{i,j} \) with three preemption points (PP_{i,1}, PP_{i,2}, and PP_{i,3}). Also, we show three preempting jobs of \( \tau_h \) such that \( P_h > P_i \).

In Fig. 2, it is shown that immediately after the NPR, with the maximum blocking \( b_i \) on \( \tau_{i,j} \), starts to execute, \( \tau_i \) and \( \tau_h \) are released. Further, we distinguish among three important time intervals. With \( S_{i,j} \) we denote the latest start time of the last NPR of \( \tau_{i,j} \), and with \( F_{i,j} \) we denote its latest finish time. However, notice that we also denote \( I_i \), which is the maximum duration from the start time of the first NPR of \( \tau_{i,j} \) until the start time of the last NPR of \( \tau_{i,j} \). Therefore, this interval also...
represents the maximum time interval during which higher priority jobs can affect CRPD of \( \tau_{i,j} \) due to preemptions. Note that the latest start time \( S_{i,j} \) of the last NPR includes lower priority blocking (\( b_{i} \) time units long), and possibly some higher priority interference prior to the start time of \( \tau_{i,j} \). During those two events, \( \tau_{i,j} \) cannot be preempted nor can its CRPD be affected. Therefore, \( I_{i} \) can be less than \( S_{i,j} \) and \( F_{i,j} \) and its precise estimation is important since by an over-approximation of \( I_{i} \) it is possible to account for the CRPD impact from multiple preemptions which in reality cannot occur.

\[
\begin{align*}
\tau_{1} & = \{1, 2, 3, 4\} \\
& \quad \text{ECB}_{i} = \{1, 2, 3, 4\} \\
& \quad \text{RRT} = 1 \\
\tau_{2} & = \{1, 3\} \\
& \quad \text{ECB}_{i} = \{1, 3, 4\} \\
\tau_{3} & = \{1, 2, 3\} \\
& \quad \text{ECB}_{i} = \{1, 2, 3, 4\}
\end{align*}
\]

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& \quad \text{ECB}_{i} = \{1, 2, 3, 4\}
\end{align*}
\]

In Fig. 3, we show a low-level perspective of the tasks, with associated task and cache usage parameters. The accessed cache-sets are depicted as circled integers in the NPRs where they are accessed. For example, task \( \tau_{1} \) accesses cache blocks in cache-sets 1, 2, and 3, and 4. Considering task \( \tau_{3} \), in its first NPR it accesses cache blocks in cache sets 1 and 2, and so on. Therefore, a set of useful cache blocks at \( PP_{3,1} \) is \( UCB_{3,1} = \{1, 2\} \), while at \( PP_{3,2} \) is \( UCB_{3,2} = \{1, 2, 3\} \), and at \( PP_{3,3} \) it is \( UCB_{3,3} = \{1, 2, 3, 4\} \).

To motivate the necessity for the analysis on this level, assume that \( \tau_{1} \) preempts \( \tau_{3} \) at all of its three preemption points. Let us now focus on cache-set 1. Although it is considered that there might be a ECB residing in cache-set 1 at each of the three preemptions points, this cache block can be reloaded at most once due to preemption. This reload can only occur within the last NPR of \( \tau_{3} \) since only in the last NPR it is re-accessed. To consider the worst-case eviction at each of those three points in isolation, as in method proposed by Yao et al.\[11\], can introduce over-approximation, accounting for 3 block reloads in cache set 1, opposed to at most one possible reload. Next to this, we also consider several more properties of task interactions on the cache level which are defined and formalised in the following analysis.

V. CACHE-AWARE RESPONSE TIME ANALYSIS

Considering the computation of the response time \( R_{i} \) of \( \tau_{i} \), the proposed analysis accounts for the following main computations (see Fig. 2):

1) \( b_{i} \) – Computation of the maximum lower-priority blocking on \( \tau_{i} \), accounting for the CRPD in the blocking.
2) \( I_{i} \) – Computation of the maximum time duration from the start time of \( \delta_{i,j} \) until the start time of \( \delta_{i,j} \). \( I_{i} \) represents the upper bound on the interval during which higher priority jobs can affect CRPD of \( \tau_{i,j} \) due to preemptions (see Fig. 2).
3) \( S_{i,j} \) – Computation of the latest start time of the last NPR of \( \tau_{i,j} \) accounting for the CRPD on all jobs with higher priority than \( \tau_{i} \) which can be released within \( S_{i,j} \), and accounting for the CRPD on \( j \) – 1 jobs of \( \tau_{i} \). The CRPD of \( \tau_{i,j} \) should only include the CRPD in first \( l_{i} - 1 \) NPRs, since in the last NPR cache block reloads occur after \( S_{i,j} \), during the execution of \( \delta_{i,l_{i}} \). (see Fig. 2)
4) \( F_{i,j} \) – Computation of the latest finish time of the j-th job of \( \tau_{i} \), where the worst-case execution time of the last NPR \( \delta_{i,l_{i}} \) of \( \tau_{i} \) is included, with the CRPD which can be caused during the execution of \( \delta_{i,l_{i}} \).

In the remainder of this section, we formally define the terms for computing the response time of a task, divided into three subsections: 1) Maximum lower priority blocking, 2) CRPD computations, and 3) Latest start time, finish time, and response time. In order to ease the understanding of the following equations, the motivating example from Fig. 3 will also serve as the running example.

A. Maximum lower priority blocking

We start by defining the longest worst-case execution time of a NPR in \( \tau_{i} \) including the time the NPR may be prolonged due to the CRPD.

\[
\begin{align*}
\delta_{i}^{\text{max}}_{t} & = \max_{k \in [1,h] \setminus \{t\}} \left( q_{i,k} + \left| UCB_{i,k-1} \right| \cap ECB_{i,k} \right) \bigcup \left( \bigcup_{t \in k} ECB_{t} \times RRT \right), \quad \text{where } UCB_{i,0} = \emptyset
\end{align*}
\]

\( \delta_{i}^{\text{max}}_{t} \) is an upper bound on the longest worst-case execution time, accounting for CRPD, of a non-preemptive region in \( \tau_{i} \).
Proof: The CRPD of a non-preemptive region $\delta_{i,k}$ cannot be higher than BRT multiplied by the number of cache blocks such that: 1) the block can be evicted by some task with higher priority than $\tau_i$ at $PP_{i,k-1}$, and 2) due to preemption, the block can be reloaded in $\delta_{i,k}$, i.e. can be accessed in $\delta_{i,k}$, and is useful at $PP_{i,k-1}$. This is accounted in Equation 1 by the set intersections between $UCB_{i,k-1}$, $ECB_{i,k}$, and the union of all evicting cache blocks from the tasks with higher priority than $\tau_i$. Since the equation computes the maximum sum of the WCET and the CRPD upper bound, from all NPRs in $\tau_i$, then the proposition holds.

Example: Given the task $\tau_3$ from Fig. 3, $q_{i}^{max}=8$, since the NPR $\delta_{3,4}$ has the non-preemptive WCET equal to 4, and all of its 4 useful cache blocks might be reloaded during its execution due to preemption at $PP_{3,3}$.

Given the $q_{i}^{max}$ term, we now define the maximum lower priority blocking that may block the start time of an instance of $\tau_i$. Informally, it is the longest worst-case execution time of a NPR belonging to the one of the tasks with a lower priority than $\tau_i$.

Definition 4: Maximum lower priority blocking $b_i$ which may delay the start time of an instance of $\tau_i$ is equal to:

$$b_i = \max_{\tau_j \in \{\tau_i, \ldots, \tau_k\}} (q_{j}^{max})$$

Example: Given the task $\tau_1$, its maximum lower priority blocking is equal to 8, since the NPR $\delta_{3,4}$ has the largest WCET value accounting for possible CRPD, among all the other NPRs of $\tau_2$ and $\tau_3$.

B. CRPD related computations

This subsection is divided into four main parts:

1) Reloadable cache blocks – where a modification on the term of useful cache blocks in tasks with NPRs is defined.

2) CRPD of a single job of $\tau_i$ during the time interval $t$ – Where the computations for the upper bound on CRPD of a single job during the time interval $t$ are defined.

3) CRPD of all jobs of $\tau_i$ released during the time interval $t$ – Where the computations for the upper bound on CRPD for all jobs of a task that can be released during the time interval $t$ are defined.

4) Maximum time interval $I_{i}$ – Where the computations for the maximum time duration from the start time of $\delta_{i,1}$ until the start time of $\delta_{i,1}$ is defined.

1) Reloadable cache blocks:

In this part we first define an upper bound on the number of times a cache memory block $m$ can be reloaded during the execution of a task. We exploit the fact that although a single cache block $m$ may be useful and evicted at several consecutive preemption points, it does not necessarily mean that it can be reloaded as many times as it is useful. We prove that between the two consecutive accesses of $m$ in $\tau_i$, there is at most one cache block reload caused by preemptions between the accesses.

Lemma 1: If a memory block $m$ is not accessed in a NPR $\delta_{i,k}$, then it cannot be reloaded in $\delta_{i,k}$.

Proof: A cache block reload of $m$ in $\delta_{i,k}$ implies that it must be accessed during the execution of $\delta_{i,k}$, without being present in the cache memory, which concludes the proof.

Lemma 2: A memory block $m$ can be reloaded at most once in $\delta_{i,k}$ due to preemption.

Proof: (By contradiction) Let us assume that due to preemption $m$ can be reloaded more than once in $\delta_{i,k}$. During the execution of $\delta_{i,k}$, the first access request for the memory block $m$ can result in a reload due to preemption and eviction of $m$ prior the execution of $\delta_{i,k}$. However, the next reload due to preemption within the same region implies that $\delta_{i,k}$ was preempted. This is a contradiction since $\delta_{i,k}$ is non-preemptive.

Note that $m$ can be self-evicted during the execution of $\delta_{i,k}$ and afterwards again reloaded in the same NPR which is accounted for by the worst-case execution time of the region. However, the source of such reload is not a preemption.

Proposition 2: If a memory block $m$ is accessed in a NPR $\delta_{i,k}$ then due to preemptions it can be reloaded at most once from $PP_{i,k}$ until the end of the following NPR $\delta_{i,l}$ where it is accessed.

Proof: (By contradiction) Let us assume that $m$ is accessed in $\delta_{i,k}$, and due to preemption $m$ can be reloaded more than once from $PP_{i,k}$ until the end of $\delta_{i,l}$. This implies two possible cases: 1) $m$ is reloaded in a NPR between $\delta_{i,k}$ and $\delta_{i,l}$ where it is not accessed, which contradicts Lemma 1, or 2) due to preemptions, $m$ is reloaded more than once in $\delta_{i,l}$ which contradicts Lemma 2.

Given the Proposition 2, in order to compute the upper bound on reloads of $m$ during the execution of first $x$ NPRs of $\tau_i$, we can compute the number of first $x$ NPRs for which $m$ is in the ECB set of a NPR and in the UCB set of the first succeeding preemption point.

Definition 5: The upper bound $ub_{i,x}$ on reloads of a memory block $m$ during execution of the first $x$ NPRs of $\tau_i$ is equal to:

$$ub_{i,x} = \left\{ k \mid (1 \leq k < x \leq l_i) \land m \in ECB_{i,k} \cap UCB_{i,k} \right\}$$

where $UCB_{i,k}$ is a set of cache blocks that must be cached at $PP_{i,k}$, and may be reused until the end of the $x$-th NPR.

Proposition 3: $ub_{i,x}$ is an upper bound on the number of cache block reloads of $m$ during the execution of the first $x$ NPRs of a single job of $\tau_i$.

Proof: By Proposition 2, it holds that $m$ can be reloaded at most once between the two succeeding accesses of $m$ in two different NPRs, $\delta_{i,k}$ and $\delta_{i,l}$. Also, it holds that between $\delta_{i,k}$ and $\delta_{i,l}$, there is exactly one preemption point for which $m \in UCB_{i,k} \cap ECB_{i,k}$, and it is the first preemption point succeeding $\delta_{i,k}$. Then, by counting the number of non-preemptive regions $\delta_{i,k}$ (1 $\leq k < x$) for which $m \in ECB_{i,k} \cup UCB_{i,k}$ is a safe upper bound on the number of reloads of $m$ during the first $x$ NPRS of $\tau_i$, which is accounted by $ub_{i,x}$.  


Example: Given the task $\tau_3$ from Fig. 3, the upper bound on the number of reloads $ub_{3,4}$ of a memory block 1 during the execution of the entire task is $ub_{3,4} = 1$. Notice that memory block 1 is useful at all three preemption points and can be evicted upon preemption on any of them, however it can be reloaded at most once. Next, we define:

Definition 6: A multiset $RCB_{i,x}$, where each memory block is present the maximum number of times it can be reloaded during the execution of the first $x$ NPRs of a single job of $\tau_i$:

$$RCB_{i,x} = \bigcup_{m \in ECB_i} \left( \bigcup_{ub_{i,z}^m} \{m\} \right)$$

Proposition 4: $RCB_{i,x}$ is a multiset of all possible reloadable cache blocks during the execution of the first $x$ NPRs of a single job of $\tau_i$.

Proof: Since each accessed block $m$ (given by $m \in ECB_i$) is included $ub_{i,z}^m$ times in the multiset, then by Proposition 3 all the possible reloads for each memory block during the execution of the first $x$ NPRs of a job of $\tau_i$ are accounted.

Example: Given the task $\tau_3$ from Fig. 3, the $RCB_{3,4}$ multiset is equal to $\{1, 2, 3, 4\}$, and $RCB_{3,3} = \emptyset$. Notice that memory blocks 1, 2, 3, and 4 are not present in $RCB_{3,3}$ because their reloads cannot occur during the execution of the first 3 regions of $\tau_3$. For $\tau_2$, $RCB_{2,4} = \{1, 3, 4, 4\}$.

2) CRPD of a single job of $\tau_i$ within the time interval $t$:

For the upper bound on CRPD caused on a job of $\tau_i$ during the time interval $t$, we define two complementary types of CRPD computations. In the analysis, we compute CRPD with both approaches, and later use the lower value as the final upper bound. The proposed concepts are:

- **RCB-based upper bound** ($\gamma^{\downarrow}_i$) – which considers that all higher priority jobs, releasable during $t$, can preempt a job $\tau_{i,j}$ at all preemption points. In this case, we consider that each reloadable cache block is reloaded during the execution of a job if it can be evicted by some higher priority job released during $t$.

- **Preemption-based upper bound** ($\gamma^{\uparrow}_i$) – which considers that a preemption, from a single higher priority job $\tau_{h,j}$ can evict and cause a reload only for the useful cache blocks of the point at $\tau_{i,j}$ where the preemption occurred. In this case, for each preempting task $\tau_{h}$, whose jobs can be released $j$ times during $t$ interval, we account for $j$ highest CRPD values resulting from individual preemptions of those jobs on preemption points of $\tau_{i,j}$.

To give an informal example of the two CRPD computations, let us assume that given the taskset from Fig. 3 we compute the upper bound on CRPD caused on a single job of $\tau_3$ during a 24 units long time interval. During the interval of 24 units, at most two jobs of $\tau_1$ can be released, and at most one job of $\tau_2$ can be released.

For the **RCB-based method** it does not matter how we allocate preemptions of these jobs on preemption points, we care only about what are possibly evicting cache blocks in two jobs from $\tau_1$ and one job from $\tau_2$, and those are $\{1, 1, 1, 2, 3, 3, 4, 4, 4\}$. Also, we consider what cache blocks can be reloaded upon those evictions in single job of $\tau_3$, and those are $RCB_{3,4} = \{1, 2, 3, 4\}$. For each cache block, present in both – the RCB- and the evicting cache blocks multisets, we account that it is evicted and reloaded, meaning that CRPD at the end considers four cache block relocations.

For the **preemption-based method** the allocation of preemptions on preemption points of a task is important. First, we consider two preempting jobs of $\tau_1$ and their two possible preemptions on a job of $\tau_3$. In this case, the maximum CRPD from two preemptions is caused on $PP_3$ and $PP_2$, causing a number of reloads equal to 4 (due to evictions and reloads of blocks in $UCB_{3,3}$), and 3 (due to evictions and reloads of blocks in $UCB_{3,2}$). For the preemption from a job of $\tau_2$, the maximum CRPD occurs when $PP_3$ is preempted, since useful cache blocks 1, 2, and 3 can be evicted, and CRPD accounts for 3 cache block reloads. In total, $7 + 3 = 10$ cache block relocations are accounted for, which is the less precise upper bound than the one derived with the RCB-based method.

However, in some cases, e.g., computing CRPD for a single job of $\tau_2$, the preemption-based method provides a better result. In this case, the RCB-based method would account for 3 cache block relocations from a single job of $\tau_1$. However, the preemption-based method computes a tighter upper bound on CRPD accounting for one preemption from $\tau_1$ on $PP_{2,1}$. This is the case that produces the maximum CRPD, equal to two cache block relocations of memory blocks 1, and 3. For a single preemption at any other point, the CRPD is equal to one.

It is important to mention that in this paper, for a single job of any task, we are interested in the CRPD that can be caused:

- During the execution of the first $l_i - 1$ NPRs of a job.
- During the execution of all NPRs of a job.

This is important because when considering the time intervals (see Fig.2) for the job $\tau_{i,j}$ itself, we should only consider CRPD caused during the first $l_i - 1$ of its NPRs. This is because in $I_l$ and $S_{i,j}$ intervals, we do not consider the CRPD of the last NPR of the job. But, when considering how this job affects other jobs, we must consider CRPD in all of its NPRs.

For these reasons, many equations in this paper consider the CRPD caused during the execution of the first $x$ NPRs of a job, where $x$ is the provided parameter. We distinguish among:

- $\gamma_{i}^{\downarrow}:$ for RCB-based CRPD computations, and
- $\gamma_{i}^{\uparrow}:$ for preemption-based CRPD computations.

Based on whether or not the last NPR should be included, as discussed above, $x$ can be assigned the value $l_i - 1$ or $l_i$, respectively.

We first define the upper bound on CRPD for a single job within a time interval $t$ using RCB-based computation:

**Definition 7:** The upper bound $\gamma_{i,x}^{\downarrow}(t)$ on the CRPD caused on the first $x$ NPRs of a single job of $\tau_i$ by all preempting jobs which can be released within time a interval $t$ is defined by the following equation:

$$\gamma_{i,x}^{\downarrow}(t) = \left| RCB_{i,x} \cap \bigcup_{h \in \text{bp}(i)} \left( \bigcup_{t \in T_{h}} ECB_{h} \right) \right| \times BRT \quad (5)$$
The equation accounts for the maximum number \( \lceil t/T_h \rceil \) of releases of jobs of each higher priority task \( \tau_h \) during \( t \), and includes all of their evicting cache blocks as many times as they can be released.

**Proposition 5:** \( \gamma_{i,x}^{U}(t) \) is an upper bound on the CRPD caused on the first \( x \) NPRs of a single job of \( \tau_i \) during the interval \( t \).

**Proof:** By Proposition 4, all possibly reloadable cache blocks in a single job of \( \tau_i \) during the execution of the first \( x \) NPRs are given by \( RCB_{i,x} \). By the multiset union, Equation 5 accounts for the maximum number of evicting cache blocks of all jobs with higher priority than \( \tau_i \) that can be released during \( t \). Since by the multiset intersection Equation 5 accounts that every reloadable cache block from \( RCB_{i,x} \) indeed results in a reload if there is at least one evicting block from all possibly preempting jobs released during \( t \), then \( \gamma_{i,x}^{U}(t) \) is an upper bound on CRPD caused on the first \( x \) NPRs of a single job of \( \tau_i \) during \( t \).

**Example:** Given all the NPRs of a task \( \tau_3 \) (see Fig. 3) the upper bound on their CRPD during the 24 units long time interval is equal to:

\[
\gamma_{3,4}^{U}(24) = \left\{ 1, 2, 3, 4 \right\} \cap \left( \bigcup \left\{ 1, 2, 3, 4 \right\} \cup \bigcup \left\{ 1, 2, 3 \right\} \right) \leq \left\lceil \frac{t}{T_3} \right\rceil = \left\lceil \frac{22}{22} \right\rceil = 4
\]

Before defining the preemption-based CRPD upper bound, we first define a multiset which consists of all possible CRPD values that jobs of \( \tau_h \) can cause on the first \( x \) NPRs of \( \tau_i \).

**Definition 8:** A multiset \( CRPD_{i,x}^{h} \) of all upper bounded CRPDs that can be caused by a job of \( \tau_h \) on each of the first \( x \) NPRs of a single job of \( \tau_i \) is equal to:

\[
CRPD_{i,x}^{h} = \bigcup_{k=1}^{x} \{ ECB_h \cap UCB_{i,k-1} \times BRT \}
\]

**Proposition 6:** \( CRPD_{i,x}^{h} \) is a multiset of upper bounds of the CRPD that can be caused by a job of \( \tau_h \) on the first \( x \) NPRs of a single job of \( \tau_i \).

**Proof:** A CRPD upper bound on a single NPR \( \delta_{i,k} \) is computed by accounting that all useful cache blocks at the preemption point directly preceding \( \delta_{i,k} \) can be evicted by the evicting cache blocks of \( \tau_h \). Thus, CRPD on \( \delta_{i,k} \) is computed as the intersection between the \( ECB_h \) and \( UCB_{i,k-1} \) sets. Since Equation 6 combines the upper bounds from the first NPRs in a multiset, the proposition holds.

**Example:** Given the taskset from Fig. 3, a multiset \( CRPD_{1,4}^{3,4} \) of CRPDs from a job of \( \tau_1 \) caused on the first 4 NPRs of \( \tau_3 \) is:

\[
CRPD_{1,4}^{3,4} = \left\{ \emptyset, \left\{ 1, 2 \right\}, \left\{ 1, 2, 3 \right\}, \left\{ 1, 2, 3, 4 \right\} \right\} = \{ 0, 2, 3, 4 \}
\]

Considering the cumulative CRPD caused by preemptions from jobs of \( \tau_h \) during a time interval \( t \) we define the term:

**Definition 9:** The upper bound \( \gamma_{i,x}^{U}(t) \) on CRPD caused by preemptions from jobs of \( \tau_h \) on the first \( x \) NPRs of a single job of \( \tau_i \), within the time interval \( t \), is defined in the following equation:

\[
\gamma_{i,x}^{U}(t) = \begin{cases} \sum CRPD_{i,x}^{h}, & \text{max} \left\{ \left\{ \sum c \in CRPD_{i,x}^{h} \wedge \left\lceil t/T_h \right\rceil \right\} \right\} \\ \text{otherwise} \end{cases}
\]

**Equation 7** computes the maximum cumulative CRPD from a maximum number of preemptions of jobs of \( \tau_h \) on the first \( x \) NPRs of an instance of \( \tau_i \). Equation 7 considers two cases:

**Case 1:** \( t \leq \left\lceil \frac{1}{T_h} \right\rceil \), the number of preempting jobs of \( \tau_h \) which can cause CRPD on first \( x \) regions of a single job of \( \tau_i \) is greater than or equal to the number of preemption points that can be preempted.

**Lemma 3:** \( \gamma_{i,x}^{h}(t) \) is an upper bound on CRPD caused by preemptions from jobs of \( \tau_h \) on the first \( x \) NPRs of a single job of \( \tau_i \), within the time interval \( t \), in case \( t \leq \left\lceil \frac{1}{T_h} \right\rceil \).

**Proof:** When \( t \leq \left\lceil \frac{1}{T_h} \right\rceil \), Equation 7 sums all the CRPD values in the multiset \( CRPD_{i,x}^{h} \), and since by Proposition 6, \( CRPD_{i,x}^{h} \) is a multiset of upper bounds of the CRPD that can be caused by a job of \( \tau_h \) on the first \( x \) NPRs of a single job of \( \tau_i \), then the lemma holds.

**Example for** \( \left\lceil \frac{1}{T_3} \right\rceil < x \): Given the taskset from Fig. 3, the upper bound on CRPD caused by preemptions from jobs of \( \tau_3 \) on all NPRs of a single job of \( \tau_3 \) during 100 time units is equal to \( \gamma_{3,4}(100) = 9 \) although jobs of \( \tau_1 \) can be released 5 times during 100 time units, given by \( [100/T_1] = [100/22] \). This is the case because only three preemption points can be preempted, and each of those preemptions individually cause maximum CRPD of 2, 3, and 4 cache block reloads respectively.

**Case 2:** \( t > \left\lceil \frac{1}{T_h} \right\rceil \), the number of jobs which can cause CRPD on the first \( x \) regions of a single instance of \( \tau_i \) is less than the number of preemption points of \( \tau_i \).

**Lemma 4:** \( \gamma_{i,x}^{h}(t) \) is an upper bound on CRPD caused by preemptions from jobs of \( \tau_h \) on the first \( x \) NPRs of a single job of \( \tau_i \), within the time interval \( t \), in case \( x > \left\lceil \frac{1}{T_h} \right\rceil \).

**Proof:** By Proposition 6, \( CRPD_{i,x}^{h} \) is a multiset of upper bounds of the CRPD that can be caused by a job of \( \tau_h \) on the first \( x \) NPRs of a single job of \( \tau_i \). Since \( \left\lceil \frac{1}{T_h} \right\rceil \) is the maximum number of preemptions from the jobs of \( \tau_h \) during \( t \), then only \( \left\lceil \frac{t}{T_h} \right\rceil \) number of NPRs can experience CRPD caused by jobs of \( \tau_h \). Since Equation 7 sums the \( \left\lceil \frac{1}{T_h} \right\rceil \) highest values from \( CRPD_{i,x}^{h} \) multiset, then the lemma holds.

**Example for** \( \left\lceil \frac{1}{T_h} \right\rceil < x \): Given the taskset from Fig. 3, the upper bound on CRPD caused by preemptions from jobs of \( \tau_1 \) on all NPRs of a single job of \( \tau_3 \) during 20 time units is equal to \( \gamma_{3,4}(20) = 4 \). This is the case because only one job of \( \tau_1 \) can be released during 20 time units, given by \( [20/T_1] = [20/22] \), meaning that only one preemption point can be preempted. Therefore, Equation 7 picks the single highest value in \( CRPD_{1,4}^{3,4} \), which is 4 (representing the case when \( PP_{3,3} \) is preempted).

**Proposition 7:** \( \gamma_{i,x}^{h}(t) \) is an upper bound on CRPD caused by preemptions from jobs of \( \tau_h \) on the first \( x \) NPRs of a single job of \( \tau_i \), within the time interval \( t \).

**Proof:** Follows directly from Lemma 3, and 4.
Now, we consider that all higher priority jobs may contribute to the CRPD of a single job of $\tau_i$ during the time interval $t$.

**Definition 10:** The upper bound $\gamma_{i,x}^h(t)$ on CRPD caused by preemptions of all higher priority jobs during the time interval $t$ on the first $x$ NPRs of a single job of $\tau_i$ is defined in the following equation:

$$\gamma_{i,x}^h(t) = \sum_{\forall \tau_h \in hp(\tau_i)} \gamma_{i,x}^h(t)$$

**Proposition 8:** $\gamma_{i,x}^h(t)$ is an upper bound on CRPD caused by all higher priority jobs which can be released within time interval $t$ on the first $x$ NPRs of a single job of $\tau_i$.

**Proof:** By Proposition 7, for a single higher priority task $\tau_h$, $\gamma_{i,x}^h(t)$ is an upper bound on CRPD caused by all jobs of $\tau_h$ on the first $x$ NPRs of a single job of $\tau_i$ within $t$. Since Equation 8 sums $\gamma_{i,x}^h(t)$ for each task with higher priority than $\tau_i$, the proposition holds.

**Example:** Given the taskset from Fig. 3: $\gamma_{3,4}^+(24) = \gamma_{3,4}^1(24) + \gamma_{3,4}^2(24) = 7 + 3 = 10$. This is the case because Equation 8 accounts for two preemptions from $\tau_1$ resulting in adding the two highest CRPD values (occurring at points $PP_{3,3}$: four reloads, and $PP_{3,2}$: three reloads), and one preemption from $\tau_2$ resulting in a single highest CRPD value (occurring at $PP_{3,3}$: three reloads).

Finally, we define the CRPD upper-bound using both – the RCB-based and the preemption-based computations.

**Definition 11:** Combining the two proposed methods, a safe upper bound $\gamma_{i,x}(t)$ on CRPD caused on the first $x$ NPRs of a job of $\tau_i$ during the time interval $t$ is derived with the following equation:

$$\gamma_{i,x}(t) = \min\left(\gamma_{i,x}^+(t), \gamma_{i,x}^h(t)\right)$$

**Proposition 9:** $\gamma_{i,x}(t)$ is an upper bound on CRPD caused on the first $x$ NPRs of a job of $\tau_i$ during the time interval $t$.

**Proof:** $\gamma_{i,x}(t)$ is an upper bound because it takes the least of two upper bounds, given by Propositions 5 and 8.

**Example:** Given the taskset from Fig. 3, the final safe upper bound is $\gamma_{3,4}(24) = \min(\gamma_{3,4}^+(24), \gamma_{3,4}^+(24)) = \min(4,10) = 4$, meaning that the tighter upper bound is given by the RCB-based method. However, in case of $\tau_2$, the CRPD on a single job within 100 time units is computed as follows: $\gamma_{2,4}(10) = \min(\gamma_{2,4}^+(10), \gamma_{2,4}^h(10)) = \min(3,2) = 2$, meaning that the tighter upper bound is given by the preemption-based method.

3) **CRPD on all jobs of $\tau_i$ released within time interval $t$:**

Considering that $I_i$ is the maximum time interval during which jobs with higher priority than $\tau_i$ may affect the CRPD of a single job of $\tau_i$ (see Fig. 2), then the CRPD upper bound of all jobs of $\tau_i$ which can be released during the time interval $t_i$ is defined as follows:

**Definition 12:** The upper bound $\gamma_i(t)$ on CRPD on all the jobs of $\tau_i$ that can be released during the time interval $t$ is defined with the following equation:

$$\gamma_i(t) = \left\lceil \frac{t}{I_i} \right\rceil \times \gamma_{i,l-1}(I_i)$$

**Proposition 10:** Assuming that $I_i$ is an upper bound on the time duration from the start of the first NPR of $\tau_i$ until the start of the last NPR of $\tau_i$, then $\gamma_i(t)$ is a CRPD upper bound on all the jobs of $\tau_i$ which can be released during $t$.

**Proof:** By Proposition 9 and the assumption in Proposition 10, Equation 10 derives the CRPD upper bound $\gamma_{i,l-1}(I_i)$ caused during the $I_i$ interval on all NPRs of a single job of $\tau_i$. Since for each job of $\tau_i$, that can be released within $t$, its CRPD upper bound $\gamma_i(t)$ is multiplied by the maximum number of releases, then $\gamma_i(t)$ is a CRPD upper bound on all the jobs of $\tau_i$ which can be released during $t$.

Computation of $I_i$ is defined in Equation 11.

4) **Maximum time interval $I_i$:**

Next, we define the computation of the $I_i$ interval. In Equation 11, we consider the critical instant, i.e. that just after the start of execution of $\delta_{i,1}$ all the higher priority jobs may be released. It also considers the execution of a job of $\tau_i$ until the start of its last NPR, accounting for the possible CRPD during the execution of the first $l_i-1$ NPRs. In each new iteration, the Equation accounts that the computed interval can be enlarged with the additional CRPD by the newly released higher priority jobs with their respective CRPDs.

**Definition 13:** The upper bound $I_i$ of the time duration from the start time of $\delta_{i,1}$ until the start time of $\delta_{i,l}$ is defined as the least fixed point of the following recursive equation:

$$I_i^{(0)} = E_i$$

$$I_i^{(r)} = E_i + \gamma_{i,l-1}(I_i^{(r-1)}) + \sum_{\forall \tau_h \in hp(i)} \left(\left\lceil \frac{I_i^{(r-1)}}{T_h} \right\rceil + 1\right) \times C_h + \gamma_h(I_i^{(r-1)})$$

**Proposition 11:** $I_i$ is the upper bound of the time duration from the start of the first NPR of $\tau_i$ until the start time of the last NPR of $\tau_i$ for any job of $\tau_i$.

**Proof:** By induction over the tasks in $\Gamma$, in a decreasing priority order.

**Base case:** $I_t = E_i$, because $hp(\tau_i) = \emptyset$ and $\gamma_{i,l-1}(t) = 0$ for any $t$ since $\tau_i$ cannot experience CRPD. $E_i$ is the upper bound on the time duration from the start time of the first NPR until the start time of the last NPR of $\tau_i$ since it is the worst-case execution time between those two points.

**Inductive hypothesis:** Assume that for all $\tau_h$ in $hp(i)$, $I_h$ is an upper bound on the time duration from the start of the first NPR until the start time of the last NPR of $\tau_h$.

**Inductive step:** We show that Equation 11 computes the safe upper bound $I_i$. Consider the least fixed point of Equation 11, for which $I_i = I_i^r = I_i^{r-1}$. At this point, the equation accounts for the following upper bounds and worst-case execution times:

- $E_i$, which is the worst-case execution time, assumed by the system model, of the first $l_i - 1$ NPRs.
- $\gamma_{i,l-1}(I_i)$, which is proved by Proposition 9 to be the CRPD upper bound of a single job of $\tau_i$ in the first $l_i - 1$ NPRs during the time interval $I_i$.
- $\sum_{\forall \tau_h \in hp(i)} (\left\lceil \frac{I_i}{T_h} \right\rceil + 1) \times C_h$ – Worst-case interference
caused by the execution of jobs of higher priority tasks. The interference is accounted as the sum of individual WCET values of each job with higher priority than \( \tau_i \), that can be released during \( I_i \). The maximum number of jobs, of a single higher priority task \( \tau_h \), which can be released during \( I_i \) is accounted by term \( \lceil \frac{I_i}{T_h} \rceil + 1 \).

\[ \sum_{\forall \tau_l \in \text{hp}(i)} \gamma_l(I_i) = \text{CRPD upper bound on all jobs of the higher priority tasks that can be released during } t. \]

By Proposition 10 and the inductive hypothesis, \( \gamma_l(I_i) \) is an upper bound on the CRPD of all jobs of \( \tau_l \) which can be released within \( I_i \). Then, the sum of CRPD bounds \( \gamma_h(I_i) \) for each higher priority task than \( \tau_i \) is a safe upper bound on the cumulative CRPD in all jobs with higher priority than \( \tau_i \), released during \( I_i \).

Since we proved for all the factors, which can prolong the time duration from the start time of the first NPR until the finish time of the last NPR of \( \tau_i \), that they are accounted as the respective upper bounds in Equation 11, then their sum results in an upper bound, which concludes the proof.

Note that Equation 11 uses Equation 10, and vice versa. However, Equation 11 is always computable, because in order to compute \( I_i \), Equation 11 computes \( \gamma_h(t) \) term only for higher priority tasks than \( \tau_i \), meaning that it computes \( h(t) \) using only the values \( I_1 \) to \( I_{i-1} \). Since \( \gamma_1(t) \) always evaluates to zero for any value of \( t \), because \( \tau_1 \) cannot experience CRPD, this implies that \( I_2 \) can be computed using \( \gamma_1(t) = 0 \), which implies that \( I_3 \) can be computed using the previously computed \( I_1 = 0 \), and \( I_2 \) values, and so on until \( I_i \).

\section*{C. Latest start time, finish time, and response time}

In this subsection, we define the equations for the latest start time \( S_{i,j} \) of the last NPR a job \( \tau_{i,j} \), latest finish time \( F_{i,j} \) of \( \tau_{i,j} \), and the response time \( R_i \) of a task \( \tau_i \). All of those equations include the computation of the CRPD values for each job.

We first define the latest start time \( S_{i,j} \), the latest finish time \( F_{i,j} \) of the \( j \)-th instance of \( \tau_i \) after the critical instant.

\textbf{Definition 14:} The latest start time \( S_{i,j} \) of the last NPR of the \( j \)-th instance of \( \tau_i \) is equal to the least fixed point of the following recursion:

\begin{align*}
S_{i,j}^{(0)} &= b_i + E_i \\
S_{i,j}^{(r)} &= b_i + (j-1) \times C_i + \gamma_i((j-1) \times T_i) + E_i + \sum_{\forall \tau_{i,l} \in \text{hp}(i)} \left( \left( \frac{S_{i,j}^{(r-1)}}{T_h} \right) + 1 \right) \times C_h + \gamma_h(S_{i,j}^{(r-1)}) + \gamma_i(\lceil \frac{I_i}{T_h} \rceil + 1) \\
&+ \sum_{\forall \tau_{i,l} \in \text{hp}(i)} \left( \left( \frac{S_{i,j}^{(r-1)}}{T_h} \right) + 1 \right) \times C_h + \gamma_h(S_{i,j}^{(r-1)})
\end{align*}

In Equation 12, we consider the critical instant from Theorem 1, proved by Bril et al. [30].

\textbf{Proposition 12:} Assuming that the Level-1 pending workload is greater than zero at each time point from the critical instant until \( j \times T_i \), then starting from the critical instant, \( S_{i,j} \) is the upper bound on the latest start time of the last NPR of \( \tau_{i,j} \).

\textbf{Proof:} Based on the proposition assumption on the Level-1 pending workload, the start time of \( \tau_{i,j} \) is affected by all preceding \( j-1 \) jobs. Consider the least fixed point of Equation 12, for which \( S_{i,j} = S_{i,j}^{(r)} = S_{i,j}^{(r+1)} \). At this point, the equation accounts for the following upper bounds and worst-case execution times:

\( \diamond b_i \) - By Proposition 1 and Definition 4, \( b_i \) is an upper bound on lower priority blocking on \( \tau_i \).

\( \diamond (j-1) \times C_i \) - WCETs of the first \( j-1 \) jobs of \( \tau_i \), which execute prior to \( \tau_{i,j} \).

\( \diamond \gamma_i((j-1) \times T_i) \) - By Proposition 10, it is an upper bound on CRPD of the first \( j-1 \) jobs of \( \tau_i \).

\( \diamond E_i \), which is the worst-case execution time, assumed by the system model, of the first \( i-1 \) NPRs.

\( \diamond \gamma_{i,l-1}(I_i) \) - By Proposition 9, it is an upper bound on CRPD of a single job of \( \tau_{i,j} \) in the first \( i-1 \) NPRs during the time interval \( I_i \).

\( \diamond \sum_{\forall \tau_{i,l} \in \text{hp}(i)} \gamma_{i,l}(S_{i,j}) - \text{CRPD upper bound on all jobs of the higher priority tasks that can be released during } S_{i,j} \).

By Proposition 10, \( \gamma_{i,l}(S_{i,j}) \) is an upper bound on the CRPD of all jobs of \( \tau_l \) which can be released within \( S_{i,j} \). Then, the sum of CRPD bounds \( \gamma_h(S_{i,j}) \) for each higher priority task than \( \tau_i \) is a safe upper bound on the cumulative CRPD in all jobs with higher priority than \( \tau_i \), released during \( S_{i,j} \).

Since Equation 12 sums and considers upper bounds on all factors that can delay the start time of the last NPR of \( \tau_{i,j} \), the proposition holds.

For the computation of \( F_{i,j} \) (see Fig. 2) we need to consider the worst-case execution time of the last NPR of \( \tau_i \) including the upper bound on CRPD:

\textbf{Definition 15:} The worst-case execution time \( q_{i,l}^{\text{last}} \) of the last NPR of \( \tau_i \), including the upper bound on CRPD during its execution, is defined in the following equation:

\[ q_{i,l}^{\text{last}} = q_{i,l_1} + \max \{ UCB_{i,l_1-1} \cap \bigcup_{\forall \tau_{i,l} \in \text{hp}(i)} ECB_{i,l} \} \times BRT \]
Proposition 14: Assuming that the Level-i pending workload is greater than zero at each time point from the critical instant until $j \times T_i$, then starting from the critical instant, $F_{i,j}$ is the upper bound on the latest finish time of the NFR of $\tau_{i,j}$.

Proof: Follows from Propositions 12 and 13.

Now, we define the upper bound on the Level-i active period accounting for the CRPD.

Definition 17: The level-i active period $L_i$, accounting for the CRPD on all jobs within the period, is defined as follows:

$$
L_i^{(0)} = b_i + C_i
$$

$$
L_i^{(r)} = b_i + \sum_{\forall \tau_k \in hpe(i)} \left( \frac{(L_i^{(r-1)}/T_k)}{1} + 1 \right) \times C_k + \gamma_k(L_i^{(r-1)})
$$

(15)

Proposition 15: $L_i$ is an upper bound on Level-i period.

Proof: Similar to the proof for Proposition 12. Consider the least fixed point of Equation 15, for which $L_i = L_i^* = L_i^{-1}$. At this point, the equation accounts for the following upper bounds and worst-case execution times:

$\diamond \ b_i$ - By Proposition 1 and Definition 4, $b_i$ is an upper bound on lower priority blocking on $\tau_i$.

$\diamond \ \sum_{\forall \tau_k \in hpe(i)} \left( \frac{(L_i^{(r-1)}/T_k)}{1} + 1 \right) \times C_k$ - The worst-case execution time of all jobs of $\tau_i$ and higher priority jobs than $\tau_i$ which can be released within $L_i$. The maximum number of jobs, of a single task $\tau_k \in hpe(i)$, which can be released during $L_i$ is accounted by the following term $\frac{1}{T_k} + 1$.

$\diamond \ \sum_{\forall \tau_k \in hpe(i)} \gamma_k(L_i) - CRPD$ upper bound on all jobs of $\tau_i$ and higher priority jobs than $\tau_i$ which can be released within $L_i$. Follows from Proposition 10.

Since Equation 15 sums and considers upper bounds on all factors that can extend the duration of the Level-i active period, the proposition holds.

Finally, we can define the worst-case response time:

Definition 18: The worst case response time $R_i$ of $\tau_i$ is equal to the maximum relative latest finish time of a job $\tau_{i,j}$ which is released within $L_i$,

$$
R_i = \max_{j \in \{1, L_i/T_i\}} \{ F_{i,j} - (j - 1) \times T_i \}
$$

(16)

Theorem 2: $R_i$ is the safe upper bound on the worst-case response time of $\tau_i$.

Proof: Since each respective $F_{i,j}$ value is computed within $L_i$, then the proposition assumption from Proposition 14 holds. Then, by Theorem 1 $R_i$ is the safe upper bound on the worst-case response time of $\tau_i$, which concludes the proof.

1) Schedulability analysis: Given the above-proposed response time analysis, the taskset schedulability is determined by computing whether for each task the derived response time is less than or equal to the task deadline. I.e., if for any task $\tau_i$ in $\Gamma$ the inequality $R_i > D_i$ holds, the taskset is deemed unschedulable by the proposed analysis. As described before, in Section III-A, the schedulability test can be simplified to the first job of each task, as proposed by Yao et al. [11]. Also, the proposed approach can be adopted in another simplified schedulability test (see Theorem 3 from [11]), which is based on discontinuous points of cumulative execution-request functions, proposed by Bini and Buttazzo [31]. In this case, the cumulative execution request for $\tau_i$ and all higher priority tasks can be computed using the second line of Equation 11 by changing the term $I^{(r-1)}$ with $t$, while the maximum blocking on $\tau_i$ is computed using Equation 2.

D. Set-associative LRU caches

As shown by Altmeyer et al. [8], [9], in case of set-associative LRU caches, a cache-set may contain several useful cache blocks, e.g., $UCB_{i,2} = \{1, 2, 2, 2\}$, meaning that at $PP_{i,k}$, task $\tau_i$ contains three UCBs in cache-set 2, and one UCB in cache set 1. When a preemption point is preempted, one ECB of a preempting task is enough to evict all UCBs of the same cache-set. Therefore, multiple accesses to the same set by the preemption task do not need to appear in the ECB set, which means that the ECB set can remain the same as in direct-mapped caches. A bound on CRPD due to preemption from $\tau_i$ on $PP_{i,k}$ can be defined with: $UCB_{i,k} \cap ECB_i$, where the result is a multiset that contains each element from $UCB_{i,k}$ if it is also in $ECB_i$. Lastly, the RCB definition needs to be modified to account for each reloadable block as many times as it is present in the ECB set at each preemption point where the reload-ability condition from Equation 3 is satisfied. This is achieved by updating Equation 3 with $m \in UCB_{i,k} \cap ECB_i$.

VI. Evaluation

In this section, we present and discuss the evaluation results on the effectiveness of different approaches to identify schedulable tasksets. We primarily compare the existing feasibility analysis (Feasibility Analysis), proposed by Yao et al. [10], [11], and the proposed RTA analysis (denoted as CRPD-Fixed-PP). Additionally, we compare those two approaches with state-of-the-art CRPD analysis (Combined-Multiset) for fully-preemptive scheduling, proposed by Altmeyer et al. [8], and unsafe (NO-CRPD) which does not account for any preemption cost on preemption points, thus resulting in an unsafe approximation.

For the evaluation, we use basic cache-set configurations obtained with a low-level analysis tool, named LLVMTA [32]. LLVMTA uses LLVM machine intermediate representation after

<table>
<thead>
<tr>
<th>Program</th>
<th>ECB</th>
<th>UCB</th>
<th>Max</th>
</tr>
</thead>
<tbody>
<tr>
<td>alpcm</td>
<td>256</td>
<td>230</td>
<td>103</td>
</tr>
<tr>
<td>bs</td>
<td>43</td>
<td>23</td>
<td>20</td>
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<tr>
<td>bsort100</td>
<td>57</td>
<td>40</td>
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<td>123</td>
<td>58</td>
<td>44</td>
</tr>
<tr>
<td>compress</td>
<td>247</td>
<td>150</td>
<td>63</td>
</tr>
<tr>
<td>cover</td>
<td>256</td>
<td>38</td>
<td>15</td>
</tr>
<tr>
<td>crc</td>
<td>121</td>
<td>62</td>
<td>30</td>
</tr>
<tr>
<td>edn</td>
<td>256</td>
<td>222</td>
<td>123</td>
</tr>
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<td>47</td>
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<tr>
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<td>113</td>
<td>62</td>
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<td>15</td>
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<td>27</td>
</tr>
<tr>
<td>jfdctint</td>
<td>132</td>
<td>122</td>
<td>54</td>
</tr>
</tbody>
</table>

TABLE I: Cache configurations obtained with LLVM [32] analysis tool used on Mälardalen benchmark programs [33].
all backend compiler passes have run. The resulting machine program corresponds to a CFG reconstruction of a binary file. Cache-set configurations are obtained on real-time programs from the Mälardalen benchmark [33], given in Table I. The obtained parameters per each program are set of evicting (ECB) and definitely useful cache blocks (UCB) (in the table we show the size of each set). Also, the maximum number (Max) of definitely useful cache blocks per any preemption point of each program is obtained. The assumed direct-mapped cache memory consists of 256 sets with a line size of 8 bytes. For more details about the low-level analysis refer to [34].

In all of the experiments, we generate 2000 tasksets for each investigated parameter value. In order to include realistic cache-set usage, each task in a taskset is randomly assigned one of the distinct cache-set configurations presented in Table I. Since the task binaries were analysed individually, they all start at the same address (mapping to cache set 0). In a multi-task scheduling situation, this can hardly be the case because the ECB and UCB placement is determined by their respective locations in memory. We took this into account by randomly shifting the cache set indices, e.g. the ECB in cache set $i$ is shifted to the cache line equal to $(i + \text{random}(256)) \text{ modulo } 256$. Further task parameters were generated using the evaluation setup proposed by Altmeyer et al. [18], [8] which is used in many CRPD-related papers in order to exhaustively evaluate CRPD-based analyses. Utilisation of each task was generated using the standardised algorithm in real-time research – UUnifast, proposed by Bini and Butazzo [35]. The task periods were generated according to a uniform distribution from 5 to 500 milliseconds, representing the general spread of periods in automotive and aerospace hard real-time applications [18]. The worst-case execution time for each task was computed with the following equation $C_i = U_i \times T_i$, where $U_i$ represents the utilisation of $\tau_i$. Task deadlines were implicit, meaning that $D_i = T_i$. Task priorities were assigned according to the deadline-monotonic order. Furthermore, in order to allow for an exhaustive evaluation, we generated other parameters for LP-FPP task model and in different experiments we varied different parameters. We finish this paragraph by enlisting the default evaluation setup. The number of subjobs for each task was randomly generated according to a uniform distribution from the range $[3, 100]$ as this is a representative range of non-preemptive regions used in automotive real-time applications (where NPRs are also called runnables). The worst-case execution time of each subjob was set to $C_i/l_i$. The default taskset size was set to 6. For each preemption point $PP_{i,k}$ of a task $\tau_i$, its useful cache blocks $UCB_{i,k}$ were generated from the UCB set, assigned to $\tau_k$ from Table I. Also, we pessimistically assumed that each point of $\tau_k$ exhibits a maximum number of UCBs, given in the table. Accordingly, each evicting cache block was included at the NPRs prior and after the preemption point where it is useful. The assumed block reload time was 8 microseconds, as given in [18].

In order to increase the exhaustiveness of the performed evaluation, we used the weighted schedulability measure which also enables emulation of a 3-dimensional plot to a 2-dimensional one, as proposed by Bastoni et al. [7]. In the shown figures, we show the weighted schedulability measure $W_\rho(\rho)$, for schedulability test $y$ as a function of parameter $\rho$. For each value of $\rho$, this measure combines data for all of the tasksets generated for each utilisation level from 0.7 to 1, with a step of 0.2. For each value $\rho$ of the selected parameter, the schedulability measure is $W_\rho(\rho) = \sum_{\Gamma} (U_\Gamma \times B_\rho(\Gamma, \rho))/\sum_{\Gamma} U_\Gamma$, where $B_\rho(\Gamma, U_\Gamma)$ is a result (1 if schedulable, 0 otherwise) of a schedulability test $y$ for a taskset $\Gamma$ and parameter value $\rho$ ($U_\Gamma$ is a taskset utilisation). A method producing the highest weighted measure is the most prone to identify schedulable tasksets.

In the first experiment (Fig.4), we evaluated the impact of the taskset utilisation ($x$-axis) on the effectiveness of the analyses to determine the taskset schedulability ($y$-axis). We notice that $\text{CRPD-Fixed-PP}$ succeeds to identify significantly more schedulable tasksets compared to (Feasibility Analysis) and (Combined-Multiset), e.g. for $U_\Gamma = 0.88$, CRPD-Fixed-PP identifies 48%, Combined-Multiset 34%, and Feasibility-Analysis only 0.6% of all generated tasksets as schedulable.

In the second experiment (Fig.5), we evaluated the impact of the taskset size ($x$-axis) on the effectiveness of the analyses to determine the taskset schedulability ($y$-axis), with the weighted measure. The evaluated taskset size ranges from 3 to 10. We notice that with the increase in taskset size, all methods deteriorate in finding schedulable tasksets. However, CRPD-Fixed-PP deteriorates with the lowest rate, still being able to
find the largest number of schedulable tasksets compared to Feasibility-Analysis and Combined-Multiset.

In the third experiment (Fig.6), we evaluated the impact of the numbers of NPRs per task (x-axis) on the effectiveness of the analyses to determine the taskset schedulability (y-axis), with the weighted measure. In this experiment, we randomly generated the number of NPRs from the uniform distribution in the range $[3,x]$ where $x$ is a value on the x-axis from the figure. We notice that the greater the number of NPRs, the higher the weighted schedulability of CRPD-Fixed-PP, meaning it is more prone to identify schedulable tasksets. This is the case because the greater the number of NPRs, the lower the blocking from the lower priority tasks. In contrast, with Feasibility-Analysis it is the opposite because the greater number of NPRs, the greater estimation of preemption cost in this case.

![Fig. 6: Weighted measure at different upper bound on number of non-preemptive regions.](image)

In the fourth experiment (Fig.7), we evaluated the impact of the number of UCBs per preemption point (x-axis) on the effectiveness of the analyses to determine the taskset schedulability (y-axis), with the weighted measure. In this experiment, we removed the pessimistic consideration from the default evaluation setup that all preemption points exhibit the maximum number ($Max_i$ from Table I) of UCBs at each preemption point. Therefore, for each preemption point $PP_i,k$, we randomly generated the number of UCBs from the uniform distribution in the range $[x, Max_i]$ where $x$ is a value on the x-axis from the figure, and $Max_i$ the maximum number of UCBs of $\tau_i$, derived from the configuration assigned from Table I. In cases when $x$ value exceeded $Max_i$, $Max_i$ was assigned. We notice that only CRPD-Fixed-PP is affected in this experiment. In case when the number of UCBs per point is derived from the range $[0, Max_i]$ the proposed method is more prone to identify schedulable taskset, as it accounts for the variability of CRPD throughout the task’s execution, unlike the other presented methods.

![Fig. 7: Weighted measure at different upper bound on number of UCBs per preemption point.](image)

In the fifth experiment (Fig.8), we evaluated the impact of changing the reload-ability condition on CRPD-aware methods. We considered two cases. Bars (C.M. – Combined-Multiset) and (C.F.P. – CRPD-Fixed-PP) represent the weighted measures obtained with default evaluation setup, where each UCB within the task’s execution is definitely reloadable (since ECBs are assigned in NPRs directly preceding and succeeding the point where the block is useful). Bar (C.F.P’. – CRPD-Fixed-PP) considers the case where UCB is reloadable only once throughout the entire sequence of preemption points where it is useful, thus resulting in more identified schedulable tasksets. This is even more important for tasks with variable initialization in the beginning, and dispatching at the task’s end.

![Fig. 8: Weighted measure at different reload-ability conditions.](image)

**VII. Conclusions**

In this paper, we propose a cache-aware response time analysis for tasks with non-preemptive regions (NPR) under fixed-priority preemptive real-time scheduling. The analysis accounts for a tighter estimation of cache-related preemption delays thus dominating the existing feasibility analysis for the fixed-point–based tasks. It is based on the observations that: 1) the number of cache blocks which can be reloaded during the execution of a task can be smaller than the number of cache blocks which can be useful in a task, 2) a single job with NPRs can be affected by cache-related preemption delay from the preemptions occurring only during the maximum time duration from the start of its first NPR, until the start of its last NPR.

We evaluated the proposed analysis by comparing its effectiveness to identify schedulable tasksets with state of the art feasibility analysis for tasks with non-preemptive regions and fully-preemptive tasks. The evaluation showed that the proposed analysis significantly improves the state of the art of such schedulability analyses, since it accounts for the estimation of a tighter bound on cache-related preemption delay.
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